

## SMART CONTRACT AUDIT REPORT

for

PulsarSwap

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# 1 Introduction

Given the opportunity to review the design document and related smart contract source code of the PulsarSwap protocol, we outline in the report our systematic approach to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistencies between smart contract code and design document, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts is well designed and engineered, though it can be further improved by addressing our suggestions. This document outlines our audit results.

#### 1.1 About PulsarSwap

PulsarSwap is the implementation of Time-Weighted Average Market Maker (TWAMM) that effectively combines embedded AMM, LongTerm Orders, Order Pool, and scalable reward distribution to enable not only Uniswap-like DEXs, but also other AMMs with algorithmic trading TWAP. Compared to AMM, TWAMM reduces the price slippage associated with large trades, thus reducing trader losses. The basic information of the audited protocol is as follows:

Item	Description
Name	Pulsar
Website	https://pulsarswap.com/
Туре	Solidity Smart Contract
Platform	Solidity
Audit Method	Whitebox
Latest Audit Report	April 21, 2022

Table 1.1: Basic Information of PulsarSwap

In the following, we show the Git repository of reviewed files and the commit hash value used in this audit.

https://github.com/PulsarSwap/TWAMM-Contracts.git (27b5b3b)

And this is the commit ID after all fixes for the issues found in the audit have been checked in:

https://github.com/PulsarSwap/TWAMM-Contracts.git (8c7d701)

#### 1.2 About PeckShield

PeckShield Inc. [11] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystems by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (https://t.me/peckshield), Twitter (http://twitter.com/peckshield), or Email (contact@peckshield.com).

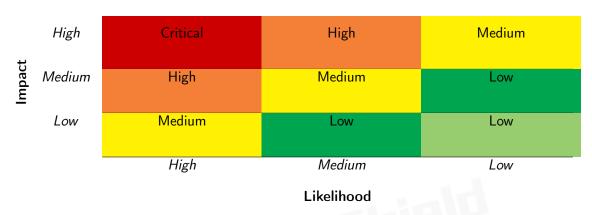


Table 1.2: Vulnerability Severity Classification

## 1.3 Methodology

To standardize the evaluation, we define the following terminology based on the OWASP Risk Rating Methodology [10]:

- <u>Likelihood</u> represents how likely a particular vulnerability is to be uncovered and exploited in the wild;
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk.

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact and can be classified into four categories accordingly, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

To evaluate the risk, we go through a checklist of items and each would be labeled with a severity category. For one check item, if our tool or analysis does not identify any issue, the contract

Table 1.3: The Full Audit Checklist

Category	Checklist Items
	Constructor Mismatch
	Ownership Takeover
	Redundant Fallback Function
	Overflows & Underflows
	Reentrancy
	Money-Giving Bug
	Blackhole
	Unauthorized Self-Destruct
Basic Coding Bugs	Revert DoS
Dasic Coung Dugs	Unchecked External Call
	Gasless Send
	Send Instead Of Transfer
	Costly Loop
	(Unsafe) Use Of Untrusted Libraries
	(Unsafe) Use Of Predictable Variables
	Transaction Ordering Dependence
	Deprecated Uses
Semantic Consistency Checks	Semantic Consistency Checks
	Business Logics Review
	Functionality Checks
	Authentication Management
	Access Control & Authorization
	Oracle Security
Advanced DeFi Scrutiny	Digital Asset Escrow
Advanced Deri Scrutilly	Kill-Switch Mechanism
	Operation Trails & Event Generation
	ERC20 Idiosyncrasies Handling
Additional Recommendations	Frontend-Contract Integration
	Deployment Consistency
	Holistic Risk Management
	Avoiding Use of Variadic Byte Array
	Using Fixed Compiler Version
	Making Visibility Level Explicit
	Making Type Inference Explicit
	Adhering To Function Declaration Strictly
	Following Other Best Practices

is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

In particular, we perform the audit according to the following procedure:

- <u>Basic Coding Bugs</u>: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- <u>Semantic Consistency Checks</u>: We then manually check the logic of implemented smart contracts and compare with the description in the white paper.
- Advanced DeFi Scrutiny: We further review business logics, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

To better describe each issue we identified, we categorize the findings with Common Weakness Enumeration (CWE-699) [9], which is a community-developed list of software weakness types to better delineate and organize weaknesses around concepts frequently encountered in software development. Though some categories used in CWE-699 may not be relevant in smart contracts, we use the CWE categories in Table 1.4 to classify our findings. Moreover, in case there is an issue that may affect an active protocol that has been deployed, the public version of this report may omit such issue, but will be amended with full details right after the affected protocol is upgraded with respective fixes.

#### 1.4 Disclaimer

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.

Table 1.4: Common Weakness Enumeration (CWE) Classifications Used in This Audit

Category	Summary
Configuration	Weaknesses in this category are typically introduced during
	the configuration of the software.
Data Processing Issues	Weaknesses in this category are typically found in functional-
	ity that processes data.
Numeric Errors	Weaknesses in this category are related to improper calcula-
	tion or conversion of numbers.
Security Features	Weaknesses in this category are concerned with topics like
	authentication, access control, confidentiality, cryptography,
	and privilege management. (Software security is not security
	software.)
Time and State	Weaknesses in this category are related to the improper man-
	agement of time and state in an environment that supports
	simultaneous or near-simultaneous computation by multiple
5 C IV	systems, processes, or threads.
Error Conditions,	Weaknesses in this category include weaknesses that occur if
Return Values,	a function does not generate the correct return/status code,
Status Codes	or if the application does not handle all possible return/status
Describes Management	codes that could be generated by a function.
Resource Management	Weaknesses in this category are related to improper manage-
Behavioral Issues	ment of system resources.
Denavioral issues	Weaknesses in this category are related to unexpected behaviors from code that an application uses.
Business Logic	Weaknesses in this category identify some of the underlying
Dusilless Logic	problems that commonly allow attackers to manipulate the
	business logic of an application. Errors in business logic can
	be devastating to an entire application.
Initialization and Cleanup	Weaknesses in this category occur in behaviors that are used
mitialization and Cicanap	for initialization and breakdown.
Arguments and Parameters	Weaknesses in this category are related to improper use of
Barrieros aria i aramieses	arguments or parameters within function calls.
Expression Issues	Weaknesses in this category are related to incorrectly written
,	expressions within code.
Coding Practices	Weaknesses in this category are related to coding practices
3	that are deemed unsafe and increase the chances that an ex-
	ploitable vulnerability will be present in the application. They
	may not directly introduce a vulnerability, but indicate the
	product has not been carefully developed or maintained.

# 2 | Findings

#### 2.1 Summary

Here is a summary of our findings after analyzing the implementation of the PulsarSwap smart contracts. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logic, examine system operations, and place DeFi-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings
Critical	0
High	1
Medium	3
Low	2
Informational	0
Total	6

We have so far identified a list of potential issues: some of them involve subtle corner cases that might not be previously thought of, while others refer to unusual interactions among multiple contracts. For each uncovered issue, we have therefore developed test cases for reasoning, reproduction, and/or verification. After further analysis and internal discussion, we determined a few issues of varying severities need to be brought up and paid more attention to, which are categorized in the above table. More information can be found in the next subsection, and the detailed discussions of each of them are in Section 3.

#### 2.2 Key Findings

Overall, these smart contracts are well-designed and engineered, though the implementation can be improved by resolving the identified issues (shown in Table 2.1), including 1 high-severity vulnerability, 3 medium-severity vulnerabilities, and 2 low-severity vulnerabilities.

Title ID Severity Category **Status** PVE-001 High Missing Access Control in Pair Security Features Resolved **PVE-002** Resolved Low Accommodation of Non-ERC20-**Coding Practices** Compliant Tokens **PVE-003** Medium Improved WETH Token Handling in Business Logic Resolved **TWAMM PVE-004** Time and State Confirmed Medium Out-of-Gas Risk In executeVirtualOrdersUntilCurrentBlock() Medium **PVE-005** Lack Of Calling updatePrice() In **Business Logic** Resolved Pair::provideInitialLiquidity() Incorrect orderldStatusMap Update Business Logic Resolved **PVE-006** Low Logic In withdrawProceedsFrom-LongTermSwap()

Table 2.1: Key PulsarSwap Audit Findings

Beside the identified issues, we emphasize that for any user-facing applications and services, it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms should kick in at the very moment when the contracts are being deployed on mainnet. Please refer to Section 3 for details.

## 3 Detailed Results

#### 3.1 Missing Access Control in Pair

• ID: PVE-001

• Severity: High

Likelihood: High

• Impact: High

• Target: Pair

• Category: Security Features [5]

• CWE subcategory: CWE-287 [2]

#### Description

In the PulsarSwap protocol, there is a Pair contract which implements the actual pool for any two ERC20 tokens. In the Pair contract, it provides a series of interfaces for LPs to add/remove liquidity, and for traders to exchange tokens. While examining these interfaces, we notice the existence of missed access control authorization that need to be corrected.

To elaborate, we show below the code snippet of the removeLiquidity() routine. As the name indicates, this routine is designed to remove liquidity from current pool for the given LP (identified by the to argument). It comes to our attention that there is no access control restriction enforced on this routine, which makes the removeLiquidity() routine opened to the public. As a result, anyone could invoke it to remove liquidity from the pool on behalf of any LP.

```
166
         function removeLiquidity (address to, uint256 lpTokenAmount)
167
              external
168
             override
             lock
169
170
             nonReentrant
171
172
             require(
173
                  lpTokenAmount <= totalSupply(),</pre>
174
                  "Not Enough Lp Tokens Available"
175
176
             updatePrice(reserveMap[tokenA], reserveMap[tokenB]);
178
             //execute virtual orders
```

```
179
             longTermOrders.executeVirtualOrdersUntilCurrentBlock(reserveMap);
181
             //the ratio between the number of underlying tokens and the number of lp tokens
                 must remain invariant after burn
182
             uint256 amountAOut = (reserveMap[tokenA] * lpTokenAmount) /
183
                 totalSupply();
184
             uint256 amountBOut = (reserveMap[tokenB] * lpTokenAmount) /
185
                 totalSupply();
187
             reserveMap[tokenA] = amountAOut;
188
             reserveMap[tokenB] -= amountBOut;
190
             burn(to, IpTokenAmount);
192
             IERC20(tokenA).transfer(to, amountAOut);
193
             IERC20(tokenB).transfer(to, amountBOut);
195
             emit LiquidityRemoved(to, IpTokenAmount);
196
```

Listing 3.1: Pair :: removeLiquidity()

Our further study shows that the access control authorization could be granted to the TWAMM contract which is a dedicated router for the Pair contract. Note there are some other routines share the same issue in the Pair contract. Such as the provideInitialLiquidity()/provideLiquidity()/instantSwapFromAToB()/longTermSwapFromAToB() routines, etc.

**Recommendation** Add the necessary access control authorization to the above mentioned routines in the Pair contract.

Status This issue has been fixed in this commit: 8c7d701.

## 3.2 Accommodation Of Non-ERC20-Compliant Tokens

• ID: PVE-002

Severity: Low

Likelihood: Low

• Impact: Low

• Target: Pair

• Category: Coding Practices [6]

• CWE subcategory: CWE-1109 [1]

#### Description

Though there is a standardized ERC-20 specification, many token contracts may not strictly follow the specification or have additional functionalities beyond the specification. In this section, we examine the transferFrom() routine and possible idiosyncrasies from current widely-used token contracts.

In particular, we use the popular token, i.e., ZRX, as our example. We show the related code snippet below. On its entry of transferFrom(), there is a check, i.e., if (balances[\_from] >= \_value && allowed[\_from][msg.sender] >= \_value && balances[\_to] + \_value >= balances[\_to]). If the check fails, it returns false. However, the transaction still proceeds successfully without being reverted. This is not compliant with the ERC20 standard and may cause issues if not handled properly. Specifically, the ERC20 standard specifies the following: "Transfers \_ value amount of tokens from address \_ from to address \_ to, and MUST fire the Transfer event. The function SHOULD throw unless the \_ from account has deliberately authorized the sender of the message via some mechanism."

```
64
       function transfer(address _to, uint _value) returns (bool) {
65
            //Default assumes totalSupply can't be over max (2^256 - 1).
66
            if (balances[msg.sender] >= _value && balances[_to] + _value >= balances[_to]) {
                balances[msg.sender] -= _value;
67
68
                balances[_to] += _value;
69
                Transfer(msg.sender, _to, _value);
70
                return true;
71
           } else { return false; }
72
       }
73
       function transferFrom(address _from, address _to, uint _value) returns (bool) {
74
            if (balances[_from] >= _value && allowed[_from][msg.sender] >= _value &&
                balances[_to] + _value >= balances[_to]) {
75
                balances[_to] += _value;
                balances[_from] -= _value;
76
77
                allowed[_from][msg.sender] -= _value;
78
                Transfer(_from, _to, _value);
79
                return true;
80
           } else { return false; }
```

Listing 3.2: ZRX.sol

Because of that, a normal call to transferFrom() is suggested to use the safe version, i.e., safeTransferFrom(), In essence, it is a wrapper around ERC20 operations that may either throw on failure or return false without reverts. Moreover, the safe version also supports tokens that return no value (and instead revert or throw on failure). Note that non-reverting calls are assumed to be successful. Similarly, there is a safe version of transfer() as well, i.e., safeTransfer().

In the following, we show the Pair::provideLiquidity() routine. If the ZRX token is supported as either of the tokenA/tokenB in the pool, the unsafe version of IERC20(tokenA).transferFrom(to, address(this), amountAIn); (line 158 and 159) may return false if the pool (spender) does not have enough allowance from the token owner (given by the to argument). But the Pair::provideLiquidity () routine expects the transferFrom() to revert on failure. Based on this, we may intend to replace the transferFrom() (line 158 and 159) with safeTransferFrom().

```
function provideLiquidity(address to, uint256 lpTokenAmount)
external
override
```

```
135
             lock
136
             nonReentrant
137
138
             require(
139
                 totalSupply() != 0,
                 "No Liquidity Has Been Provided Yet, Need To Call provideInitialLiquidity()"
140
141
             );
142
             updatePrice(reserveMap[tokenA], reserveMap[tokenB]);
143
144
             //execute virtual orders
145
             longTermOrders.executeVirtualOrdersUntilCurrentBlock(reserveMap);
146
147
             //the ratio between the number of underlying tokens and the number of lp tokens
                must remain invariant after mint
148
             uint256 amountAIn = (lpTokenAmount * reserveMap[tokenA]) /
149
                 totalSupply();
150
             uint256 amountBIn = (lpTokenAmount * reserveMap[tokenB]) /
151
                totalSupply();
152
153
             reserveMap[tokenA] += amountAIn;
154
             reserveMap[tokenB] += amountBIn;
155
156
             _mint(to, lpTokenAmount);
157
158
             IERC20(tokenA).transferFrom(to, address(this), amountAIn);
159
             IERC20(tokenB).transferFrom(to, address(this), amountBIn);
160
161
             emit LiquidityProvided(to, lpTokenAmount);
162
```

Listing 3.3: Pair::provideLiquidity()

Note the same issue also exists in other routines, such as the provideInitialLiquidity()/removeLiquidity ()/performInstantSwap() routines in the Pair contract, and the performLongTermSwap()/cancelLongTermSwap ()/withdrawProceedsFromLongTermSwap() routines in the LongTermOrders library.

**Recommendation** Accommodate the above-mentioned idiosyncrasy with safe-version implementation of ERC20-related transfer() and transferFrom().

Status This issue has been fixed in this commit: 8c7d701.

#### 3.3 Improved WETH Token Handling in TWAMM

• ID: PVE-003

Severity: MediumLikelihood: Medium

• Impact: Medium

Target: TWAMM

• Category: Business Logic [7]

• CWE subcategory: CWE-841 [4]

#### Description

As mentioned earlier, the Pair contract supports users to create pair for any two ERC20 tokens, including the WETH token which is the ERC20 wrapper for the native ETH token. To make it convenient for users to exchange with the native ETH directly, the TWAMM contract is designed as a router to take the native ETH as the input/output token. While examining these ETH related routines in the TWAMM contract, we notice the existence of improper ETH/WETH handling that need to be corrected.

To elaborate, we show below the code snippet of the TWAMM::addInitialLiquidityETH() routine. As the name indicates, this routine is designed to add initial liquidity to ETH related pool. It comes to our attention that the routine receives ETH from the LP and deposits the ETH to the WETH contract (line 101). However, the new WETH tokens are minted to the TWAMM contract, while not the LP. Moreover, after the liquidity is added, the TWAMM contract will refund the remaining ETH back to the LP (line 105). But as mentioned earlier, all the received ETH has been deposited to the WETH contract (line 101).

```
86
         function addInitialLiquidityETH(
87
             address token,
88
             uint256 amountToken,
89
             uint256 amountETH,
90
             uint256 deadline
91
         ) external payable virtual override ensure(deadline) {
92
                 IFactory(factory).getPair(token, WETH) != address(0),
93
                 "No Existing Pair Found, Create Pair First!"
94
95
96
             address pair = Library.pairFor(factory, token, WETH);
97
             (address tokenA, ) = Library.sortTokens(token, WETH);
98
             (uint256 amountA, uint256 amountB) = tokenA == token
99
                 ? (amountToken, amountETH)
100
                 : (amountETH, amountToken);
101
             IWETH10(WETH) . deposit { value : msg . value }();
102
             IPair(pair).provideInitialLiquidity(msg.sender, amountA, amountB);
103
             // refund dust eth, if any
104
             if (msg.value > amountETH) {
105
                 TransferHelper.safeTransferETH(msg.sender, msg.value - amountETH);
106
107
```

Listing 3.4: TWAMM::addInitialLiquidityETH()

What is more, when we further look into the code of the Pair::provideInitialLiquidity() routine (as shown below), we notice that the input tokens for liquidity providing are directly transferred from the LP (given by the to argument), while not the msg.sender (TWAMM in our example). That is to say, the LP shall keep both the input tokens in its balance before adding the liquidity. So in the TWAMM::addInitialLiquidityETH() routine, the WETH shall be minted to the LP, not the TWAMM.

```
86
         function provideInitialLiquidity(
87
             address to,
88
             uint256 amountA,
89
             uint256 amountB
90
         ) external override lock nonReentrant {
91
             require (
92
                 totalSupply() = 0,
93
                 "Liquidity Has Already Been Provided, Need To Call provideLiquidity()");
95
             reserveMap[tokenA] = amountA;
96
             reserveMap[tokenB] = amountB;
98
             //initial LP amount is the geometric mean of supplied tokens
99
             uint256 lpAmount = amountA
100
                 .fromUint()
101
                 . sqrt ()
102
                 .mul(amountB.fromUint().sqrt())
103
                 .toUint(); // - MINIMUM_LIQUIDITY;
104
             // _mint(address(0), MINIMUM_LIQUIDITY); // permanently lock the first
                 MINIMUM_LIQUIDITY tokens // TODO: uncomment
105
              mint(to, IpAmount);
106
             IERC20(tokenA).transferFrom(to, address(this), amountA);
107
             IERC20(tokenB).transferFrom(to, address(this), amountB);
109
             emit InitialLiquidityProvided(to, amountA, amountB);
110
```

Listing 3.5: Pair :: provideInitialLiquidity ()

Note similar issues exist in all other ETH related routines in the TWAMM contract, such as the addLiquidityETH()/withdrawLiquidityETH()/instantSwapTokenToETH()/instantSwapETHToToken() routines, etc.

**Recommendation** Revise all the ETH related routines to properly route the ETH and the WETH between the LP and the pool.

Status This issue has been fixed in this commit: 8c7d701.

## 3.4 Out-of-Gas Risk In executeVirtualOrdersUntilCurrentBlock()

• ID: PVE-004

• Severity: Medium

Likelihood: Low

• Impact: High

• Target: LongTermOrdersLib

• Category: Time and State [8]

• CWE subcategory: CWE-682 [3]

#### Description

In the PulsarSwap protocol, the executeVirtualOrdersUntilCurrentBlock() function will be triggered when a user provides liquidity, removes liquidity, performs instant swap, performs long term swap, cancels long term swap, or withdraws proceeds from long term swap. This executeVirtualOrdersUntilCurrentBlock () function will execute all pending virtual orders until current block is reached. However, if the operations that can trigger the execution of the executeVirtualOrdersUntilCurrentBlock() function does not occur for a long time, then the number of pending virtual orders could be large enough such that the subsequent execution of executeVirtualOrdersUntilCurrentBlock() could lead to out-of-gas (lines 284-291).

```
275
         ///@notice executes all virtual orders until current block is reached.
276
         function executeVirtualOrdersUntilCurrentBlock(
277
             LongTermOrders storage self,
             mapping(address => uint256) storage reserveMap
278
279
         ) internal {
280
             uint256 nextExpiryBlock = self.lastVirtualOrderBlock -
281
                 (self.lastVirtualOrderBlock % self.orderBlockInterval) +
282
                 self.orderBlockInterval;
283
             //iterate through blocks eligible for order expires, moving state forward
284
             while (nextExpiryBlock < block.number) {</pre>
285
                 executeVirtualTradesAndOrderExpiries(
286
                     self,
287
                     reserveMap,
288
                     nextExpiryBlock
289
                 );
290
                 nextExpiryBlock += self.orderBlockInterval;
             }
291
292
             //finally, move state to current block if necessary
293
             if (self.lastVirtualOrderBlock != block.number) {
294
                 executeVirtualTradesAndOrderExpiries(
295
                     self.
296
                     reserveMap,
297
                     block.number
298
                 ):
299
```

```
300 }
```

Listing 3.6: LongTermOrdersLib::executeVirtualOrdersUntilCurrentBlock()

**Recommendation** Take into consideration the scenario where there may exist a large number of virtual orders waiting to be executed.

Status This issue has been confirmed.

# 3.5 Lack Of Calling updatePrice() In Pair::provideInitialLiquidity()

ID: PVE-005

Severity: MediumLikelihood: High

• Impact: Low

• Target: Pair

• Category: Business Logic [7]

• CWE subcategory: CWE-841 [4]

#### Description

In the PulsarSwap protocol, the Pair contract provides an external provideInitialLiquidity() for users to provide initial liquidity. This function can be called by a user only when the totalSupply of the Pair contract is equal to 0. When analyzing this initial liquidity-providing routine provideInitialLiquidity (), we notice there is a lack of invoking updatePrice() to update the price accumulators before transferring assets to the contracts.

```
103
         ///@notice provide initial liquidity to the amm. This sets the relative price
             between tokens
104
         function provideInitialLiquidity(
105
             address to.
106
             uint256 amountA,
107
             uint256 amountB
108
         ) external override lock nonReentrant {
109
             require(
110
                 totalSupply() == 0,
111
                 "Liquidity Has Already Been Provided, Need To Call provideLiquidity()"
112
             );
113
             reserveMap[tokenA] = amountA;
114
115
             reserveMap[tokenB] = amountB;
116
117
             //initial LP amount is the geometric mean of supplied tokens
118
             uint256 lpAmount = amountA
119
                 .fromUint()
120
                 .sqrt()
```

```
121
                 .mul(amountB.fromUint().sqrt())
122
                 .toUint(); // - MINIMUM_LIQUIDITY;
             // _mint(address(0), MINIMUM_LIQUIDITY); // permanently lock the first
123
                 MINIMUM_LIQUIDITY tokens // TODO: uncomment
124
             _mint(to, lpAmount);
125
             IERC20(tokenA).transferFrom(to, address(this), amountA);
126
             IERC20(tokenB).transferFrom(to, address(this), amountB);
127
128
             emit InitialLiquidityProvided(to, amountA, amountB);
129
```

Listing 3.7: Pair::provideInitialLiquidity()

If the call to updatePrice() is not invoked in the provideInitialLiquidity() routine, the value of the state variable blockTimestampLast will remain 0. The calculation of timeElapsed in the subsequent call of updatePrice() will be not correct (line 85), thus the calculations for state variables priceACumulativeLast/priceBCumulativeLast will also be not correct.

```
82
        // update price accumulators, on the first call per block
83
        function updatePrice(uint256 reserveA, uint256 reserveB) private {
84
             uint32 blockTimestamp = uint32(block.timestamp % 2**32);
85
             uint32 timeElapsed = blockTimestamp - blockTimestampLast; // overflow is desired
86
             if (timeElapsed > 0 && reserveA != 0 && reserveB != 0) {
87
                 // * never overflows, and + overflow is desired
88
                 priceACumulativeLast +=
89
                     uint256(
90
                         UQ112x112.encode(uint112(reserveB)).uqdiv(uint112(reserveA))
91
                     ) *
92
                     timeElapsed;
93
                 priceBCumulativeLast +=
94
                     uint256(
95
                         UQ112x112.encode(uint112(reserveA)).uqdiv(uint112(reserveB))
96
                     ) *
97
                     timeElapsed;
            }
98
99
             blockTimestampLast = blockTimestamp;
100
             emit UpdatePrice(reserveA, reserveB);
101
```

Listing 3.8: Pair::updatePrice()

Note similar issue also exists in the executeVirtualOrders() routine of the same contract.

Recommendation Timely invoke updatePrice() for the above-mentioned functions.

Status This issue has been fixed in this commit: 8c7d701.

# 3.6 Incorrect orderIdStatusMap Update Logic In withdrawProceedsFromLongTermSwap()

ID: PVE-006Severity: Low

Likelihood: Low

Impact: Low

• Target: LongTermOrdersLib

• Category: Business Logic [7]

• CWE subcategory: CWE-841 [4]

#### Description

In the PulsarSwap protocol, a user can withdraw proceeds from his/her long term swap order and this can be done before or after the order has expired. If the order has expired, the total proceeds of this order will be sent to the user. If the order has expired, the total proceeds collected for this order will be sent to the user. If the order has not yet expired, the proceeds accumulated so far for this order will be sent to the user. When analyzing this withdrawProceedsFromLongTermSwap() routine, we notice the current update logic for mapping type orderIdStatusMap is not correct.

To elaborate, we show below its code snippet. Specifically, an order status should be updated to false only when this order has expired, instead of updating the order status to false regardless of whether the order has expired or not (line 216).

```
///@notice withdraw proceeds from a long term swap (can be expired or ongoing)
189
190
        function withdrawProceedsFromLongTermSwap(
191
            LongTermOrders storage self,
192
             address sender,
193
            uint256 orderId,
194
            mapping(address => uint256) storage reserveMap
195
        ) internal {
196
            //update virtual order state
197
             executeVirtualOrdersUntilCurrentBlock(self, reserveMap);
198
199
            Order storage order = self.orderMap[orderId];
200
             require(order.owner == sender, "Sender Must Be Order Owner");
201
202
             OrderPoolLib.OrderPool storage OrderPool = self.OrderPoolMap[
203
                 order.sellTokenId
204
            ];
205
             uint256 proceeds = OrderPool.withdrawProceeds(orderId);
206
207
             //charge LP fee
208
             uint256 proceedsMinusFee = (proceeds * (10000 - LP_FEE)) / 10000;
209
210
             require(proceedsMinusFee > 0, "No Proceeds To Withdraw");
211
             //transfer to owner
212
             IERC20(order.buyTokenId).transfer(sender, proceedsMinusFee);
213
```

```
// delete orderId from account list
// removeOrderId(self, orderId, msg.sender);
self.orderIdStatusMap[orderId] = false;
}
```

Listing 3.9: LongTermOrdersLib::withdrawProceedsFromLongTermSwap()

Recommendation Update the order status to false only when this order has expired.

Status This issue has been fixed in this commit: 8c7d701.



# 4 Conclusion

In this audit, we have analyzed the PulsarSwap design and implementation. PulsarSwap is the implementation of Time-Weighted Average Market Maker (TWAMM) that effectively combines embedded AMM, LongTerm Orders, Order Pool, and scalable reward distribution to enable not only Uniswap-like DEXs, but also other AMMs with algorithmic trading TWAP. The current code base is well structured and neatly organized. Those identified issues are promptly confirmed and addressed.

Moreover, we need to emphasize that Solidity-based smart contracts as a whole are still in an early, but exciting stage of development. To improve this report, we greatly appreciate any constructive feedbacks or suggestions, on our methodology, audit findings, or potential gaps in scope/coverage.



# References

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